

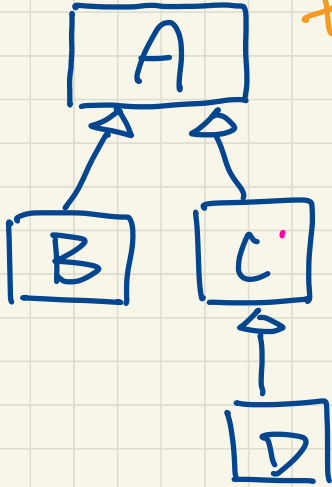
Lecture 14 - Wednesday, March 1

Announcements

- **Makeup Lecture** for WrittenTest1
 - + Expected to complete by: March 20
- **A2 solution:** only source code (no solution videos)

Static Type

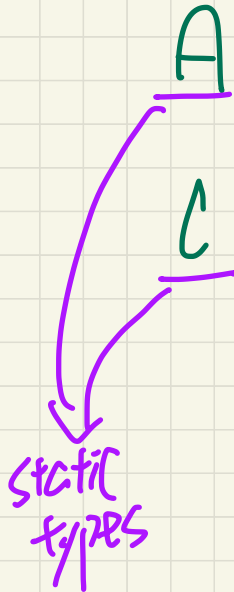
deduced type



vs. Dynamic Type

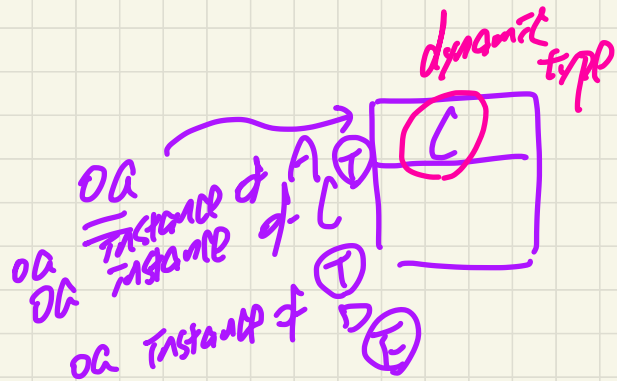
any object of type class of A? ^{dependant} any dependant classes of ST, A itself

`oa = new C();`

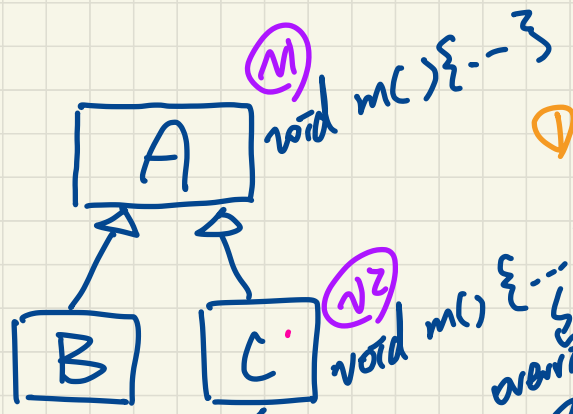


polymorphism

`oc = new D();`



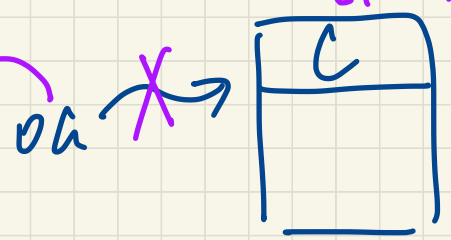
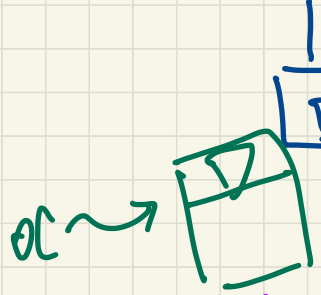
polymorphism



```

    A OA = new C();
    C OC = new D();
  
```

Annotations: (1) A, (2) C, (3) OC = new D();



```

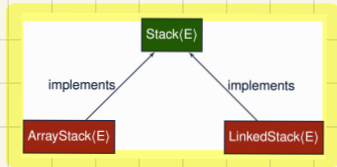
    OA.m();
  
```

Annotations: (4) OA.m();

invoke version of m in C (dynamic type)

dynamic binding
 OA.m();
 invoke version in D

Stack ADT: Testing Alternative Implementations



involve imp. in the AS push class

```

public class ArrayStack<E> implements Stack<E> {
    private final int MAX_CAPACITY = 1000;
    private E[] data;
    private t; /* index of top */
    public ArrayStack() {
        data = (E[]) new Object[MAX_CAPACITY];
        t = -1;
    }

    public int size() { return t + 1; }
    public boolean isEmpty() { return t == -1; }

    public E top() {
        if (isEmpty()) { /* Precondition Violated */ }
        else { return data[t]; }
    }

    public void push(E e) {
        if (size() == MAX_CAPACITY) { /* Precondition Violated */ }
        else { t++; data[t] = e; }
    }

    public E pop() {
        E result;
        if (isEmpty()) { /* Precondition Violated */ }
        else { result = data[t]; data[t] = null; t--; }
        return result;
    }
}
  
```

poly-morphism

involve the push imp. in LS class

```

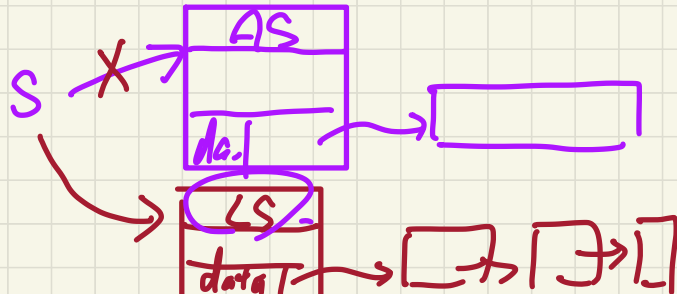
@Test
public void testPolymorphicStacks() {
    Stack<String> s = new ArrayStack<>();
    s.push("Alan"); /* dynamic binding */
    s.push("Mark"); /* dynamic binding */
    s.push("Tom"); /* dynamic binding */
    assertTrue(s.size() == 3 && !s.isEmpty());
    assertEquals("Tom", s.top());

    s = new LinkedStack<>();
    s.push("Alan"); /* dynamic binding */
    s.push("Mark"); /* dynamic binding */
    s.push("Tom"); /* dynamic binding */
    assertTrue(s.size() == 3 && !s.isEmpty());
    assertEquals("Tom", s.top());
}
  
```

ST DTI

Stack<String> s = new ArrayStack<>();

s = new LinkedStack<>();



Polymorphic Collection (stack)

```

public interface Stack<E> {
    public int size();
    public boolean isEmpty();
    public E top();
    public void push(E e);
    public E pop();
}

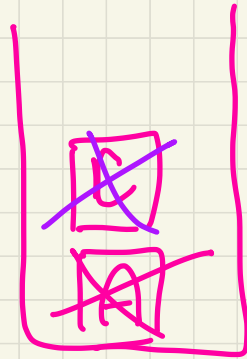
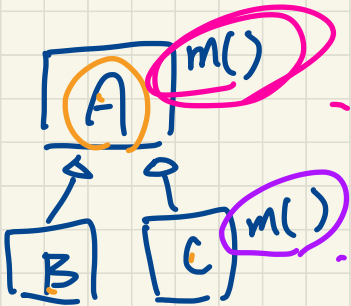
```

Stack<A> S = new ...

* S.push(new C());

any descent classes of A

??



S.push(new A());

S.push(new C());

① A obj = S.pop();

② obj.m();

③ obj = S.pop();

④ obj.m();

Lecture

Stack ADT vs. Queue ADT

***Stack ADT -
Algorithms using the Stack ADT***

Algorithm using Stack: Reversing an Array

generic parameter declared at the method level.

```
public static <E> void reverse(E[] a) {
    Stack<E> buffer = new ArrayStack<E>();
    for (int i = 0; i < a.length; i++) {
        buffer.push(a[i]);
    }
    for (int i = 0; i < a.length; i++) {
        a[i] = buffer.pop();
    }
}
```

from L to R

names

~~"Alan"~~ ~~"Mark"~~ ~~"Tom"~~

"Tom" "Mark" "Alan"

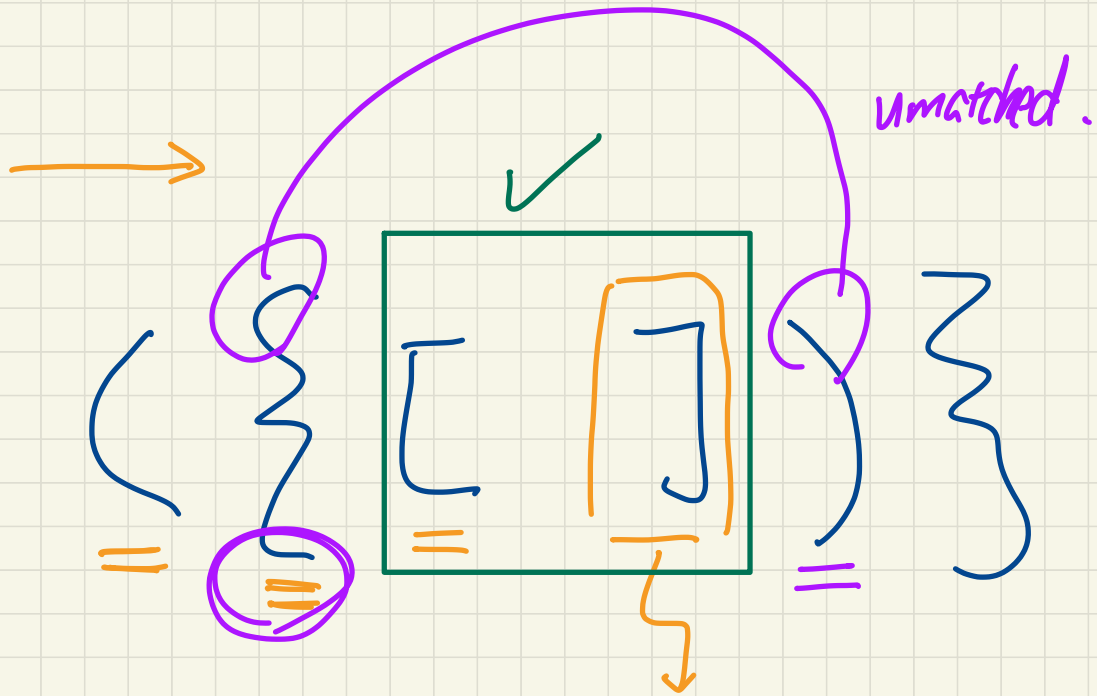
in-place reverse

~~Tom~~
~~Mark~~
~~Alan~~

buffer

```
@Test
public void testReverseViaStack() {
    String[] names = {"Alan", "Mark", "Tom"};
    String[] expectedReverseOfNames = {"Tom", "Mark", "Alan"};
    StackUtilities.reverse(names);
    assertEquals(expectedReverseOfNames, names);

    Integer[] numbers = {46, 23, 68};
    Integer[] expectedReverseOfNumbers = {68, 23, 46};
    StackUtilities.reverse(numbers);
    assertEquals(expectedReverseOfNumbers, numbers);
}
```

should match
the closest/lase
opening delimiter

Algorithm using Stack: Matching Delimiters

```

public static boolean isMatched(String expression) {
    final String opening = "[{";
    final String closing = ")]";
    Stack<Character> openings = new LinkedList<Character>();
    int i = 0;
    boolean foundError = false;
    while (!foundError && i < expression.length()) {
        char c = expression.charAt(i);
        if (opening.indexOf(c) != -1) { openings.push(c); }
        else if (closing.indexOf(c) != -1) {
            if (openings.isEmpty()) { foundError = true; }
            else {
                if (opening.indexOf(openings.top()) == closing.indexOf(c)) {
                    openings.pop();
                } else { foundError = true; }
            }
        }
        i++;
    }
    return !foundError && openings.isEmpty();
}

```

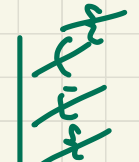
→ RT: $O(n)$
 ↓ length of input string.

closing not matched by empty stack.

↓ closing not matched

when closing matches open

openings may be non-empty



openings



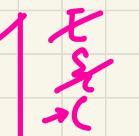
error

openings



has no matching top?

openings



end of input but stack is not empty

→ openings

```

@Test
public void testMatchingDelimiters() {
    assertTrue(StackUtilities.isMatched(""));
    assertTrue(StackUtilities.isMatched("[ ] ( )"));
    assertFalse(StackUtilities.isMatched("[ ] ( [ ] )"));
    assertFalse(StackUtilities.isMatched("[ ] ( [ ] ) ( )"));
    assertFalse(StackUtilities.isMatched("[ ] ( [ ] ) ( [ ] )"));
}

```

Post-fix notation

operands first,
then operator.

$$\underline{\underline{3}} \underline{\underline{4 \ 5 -}} * \\ 3 * (4 - 5)$$

$$\boxed{3} \boxed{4 \ 5 *} -$$

$$\underline{\underline{3 \ 4 -}} \underline{\underline{5}} * \\ (3 - 4) * 5$$

Infix notation

$$\boxed{3} \ominus \left(\underline{\underline{4}} * \underline{\underline{5}} \right)$$

operator

operands

$$\underline{\underline{5}} \underline{\underline{3 \ 4 \ominus}} * \\ 5 * (3 - 4) \\ (3 - 4) * 5$$

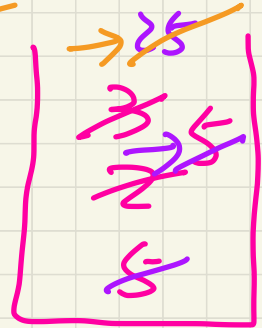
Algorithm using Stack: Calculating Postfix Expressions

Sketch of Algorithm

- When input is an **operand** (i.e., a number), **push** it to the stack.
- When input is an **operator**, obtain its two **operands** by **popping** off the stack **twice**, evaluate, then **push** the result back to stack.
- When finishing reading the input, there should be **only one** number left in the stack.

3 4 5 6 * -

? + 25

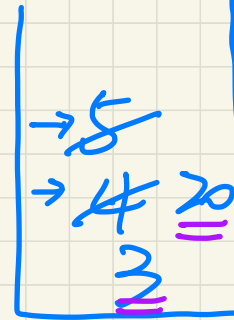


Input 1: 3 4 5 * - $\equiv 3 - (4 * 5)$

Input 2: 3 4 - 5 * $\equiv (3 - 4) * 5$

Input 3: 5 2 3 + * + $\equiv + 5 * (2 + 3)$

Input 4: 5 4 + 6 $\equiv 5 + 4 6$



$2 + 3 = 5$
 $5 * 5 = 25$

20
 $\frac{4 * 5}{LHS \quad RHS}$

$3 - 20 = -17$

Lecture

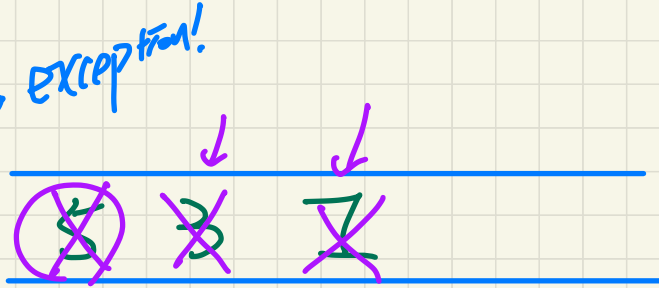
Stack ADT vs. Queue ADT

***Queue ADT -
First In First Out (FIFO)
Implementations in Java***

Queue ADT: Illustration

First-In First-Out

	isEmpty	size	first
<u>new queue</u>	T	0	
<u>enqueue(5)</u>	F	1	5
<u>enqueue(3)</u>	F	2	5
<u>enqueue(1)</u>	F	3	5
<u>dequeue</u>	F	2	3
<u>dequeue</u>	F	1	1
<u>dequeue</u>	T	0	

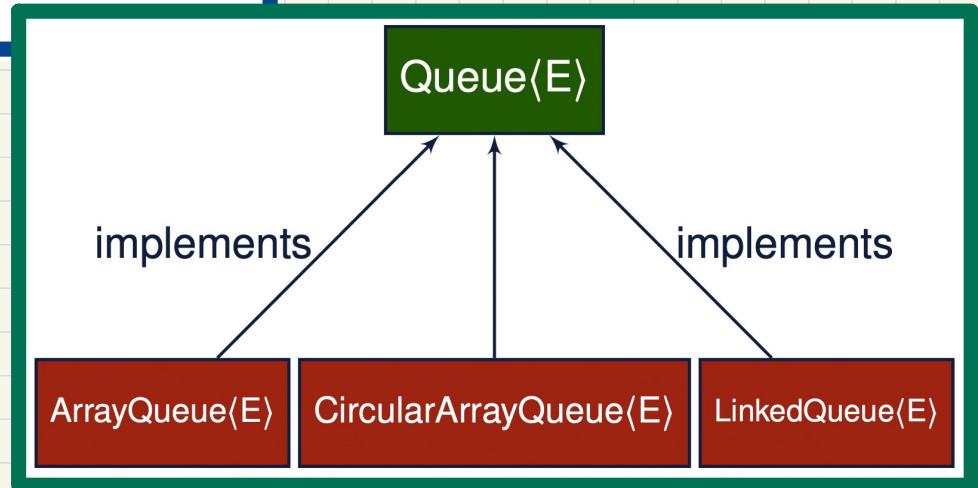


the earlier an element joins a queue, the earlier it gets removed.

exception!

Implementing the **Queue** ADT in Java: **Architecture**

```
public interface Queue< E > {  
    public int size();  
    public boolean isEmpty();  
    public E first();  
    public void enqueue( E e);  
    public E dequeue();  
}
```



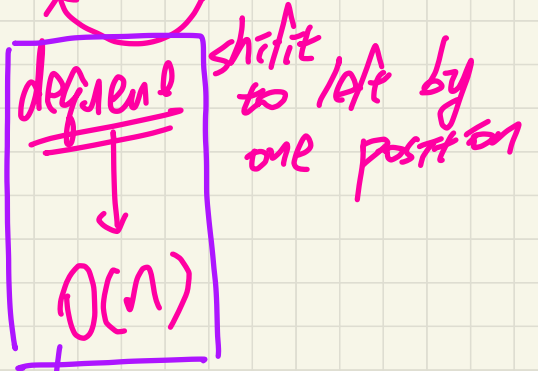
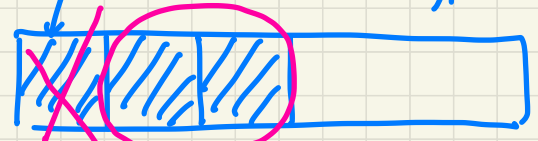
Implementing the Queue ADT using an Array

```
public class ArrayQueue<E> implements Queue<E> {
    private final int MAX_CAPACITY = 1000;
    private E[] data;
    private int r; /* rear index */
    public ArrayQueue() {
        data = (E[]) new Object[MAX_CAPACITY];
        r = -1;
    }
    public int size() { return (r + 1); }
    public boolean isEmpty() { return (r == -1); }
    public E first() {
        if (isEmpty()) { /* Precondition Violated */ }
        else { return data[0]; }
    }
    public void enqueue(E e) {
        if (size() == MAX_CAPACITY) { /* Precondition Violated */ }
        else { r++; data[r] = e; }
    }
    public E dequeue() {
        if (isEmpty()) { /* Precondition Violated */ }
        else {
            E result = data[0];
            for (int i = 0; i < r; i++) { data[i] = data[i + 1]; }
            data[r] = null; r--;
            return result;
        }
    }
}
```

$O(1)$

$O(n)$

data (beginning of array → first of Q)



to resolve this
→ circular array.
 $O(1)$